

DISJUNCTIVE PROGRAMMING.JL'S NEW SUITE OF METHODS

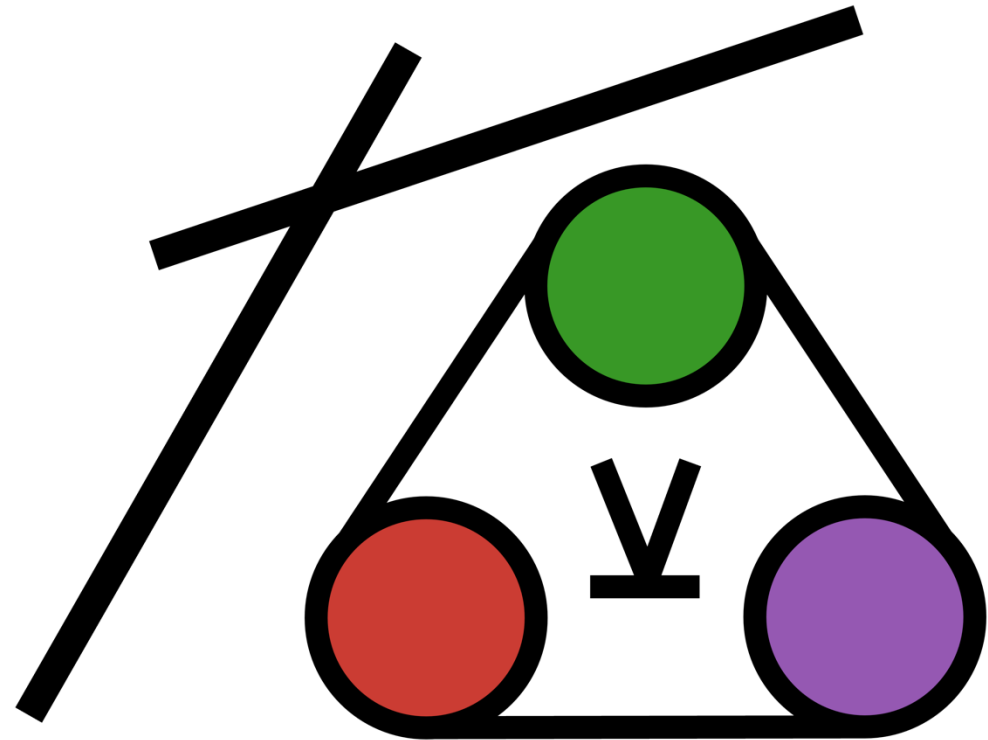
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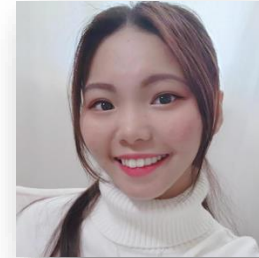
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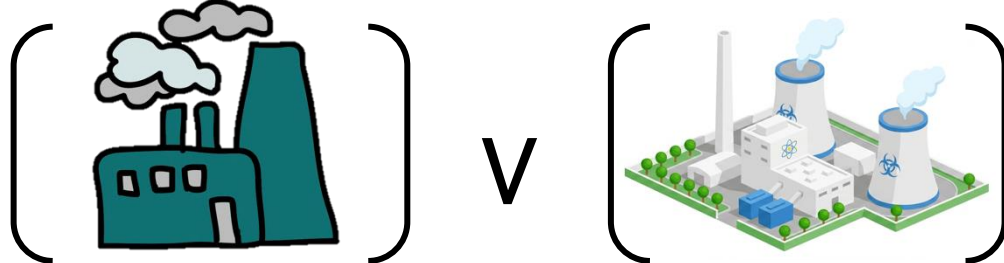


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DISJUNCTIVE PROGRAMMING OVERVIEW

Logic Constraints

- Example: Expansion Planning



- Propositional Logic:

AND	OR	NOT	IF	ONLY IF
\wedge	\vee	\neg	\Rightarrow	\Leftrightarrow

- Constraint Programming Logic:

atleast(n, \cdot) exactly(n, \cdot) atmost (n, \cdot)

- Represented via **Disjuncts** activated by **Boolean Variables**

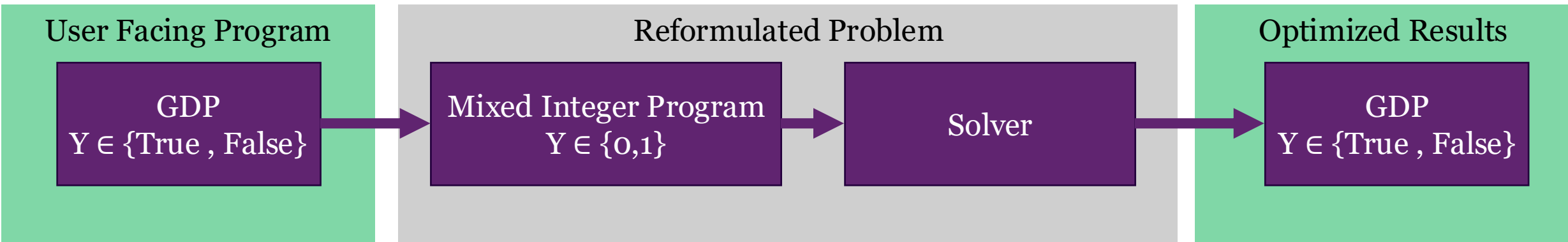
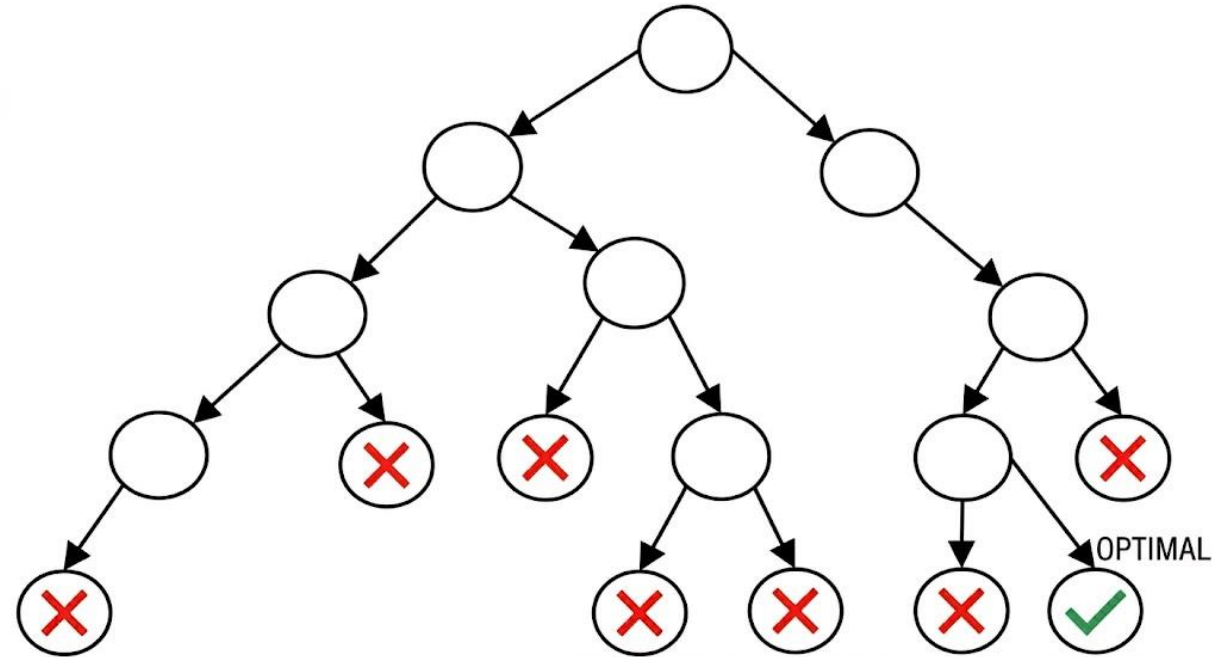
- Conditionally enforce constraints based on values of Boolean vars!**

Generalized Disjunctive Program

$$\begin{aligned}
 \min_{x, Y} \quad & f(x) \\
 \text{s.t.} \quad & g(x) \leq 0 \\
 & \bigvee_{j \in J_i} \left[r_{ij}(x) \leq 0 \right], \quad i \in \mathcal{I} \\
 & \Omega(Y) = \text{True} \\
 & x \in \mathcal{X}, \quad Y_{ij} \in \{\text{True}, \text{False}\}
 \end{aligned}$$

How to solve a GDP

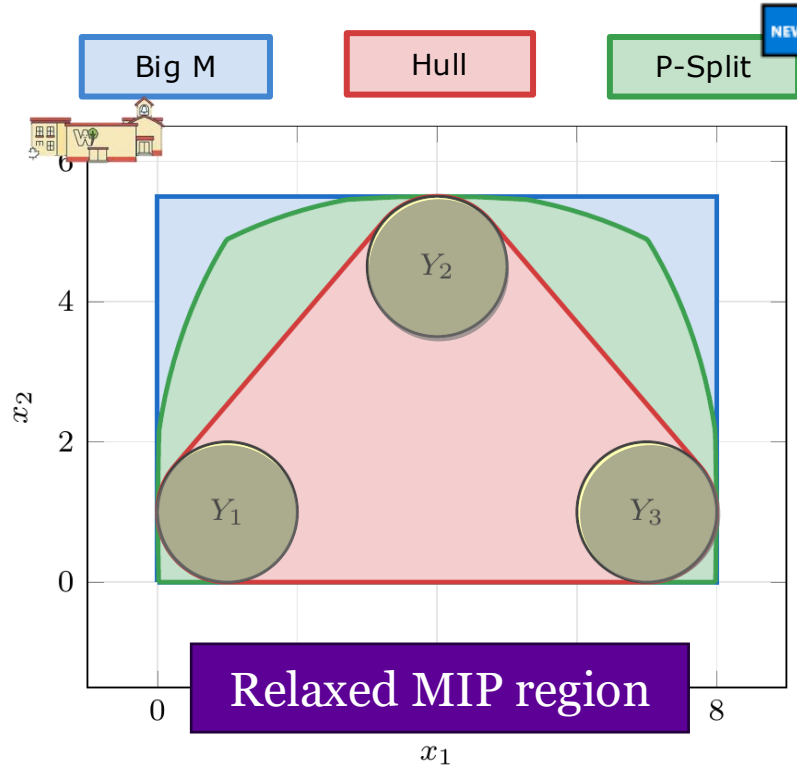
- Solvers **CANNOT** solve GDP directly
- Solvers take Mixed Integer Programs (MIPs)
- Use methods like branch and bound
 - **Integrality subject to relaxation**
- GDPs have MIP counterparts
- Recover an interpretable GDP



REFORMULATION TECHNIQUES

GDP Form

$$\begin{aligned}
 & \min_{\mathbf{x}, Y} f(x_1, x_2) \\
 & \text{s.t.} \quad \left[\begin{array}{c} Y_1 \\ (x_1-1)^2 + (x_2-1)^2 + x_3^2 + x_4^2 \leq 1 \end{array} \right] \\
 & \quad \quad \quad \vee \\
 & \quad \quad \quad \left[\begin{array}{c} Y_2 \\ (x_1-4)^2 + (x_2-4.5)^2 + x_3^2 + x_4^2 \leq 1 \end{array} \right] \\
 & \quad \quad \quad \vee \\
 & \quad \quad \quad \left[\begin{array}{c} Y_3 \\ (x_1-7)^2 + (x_2-1)^2 + x_3^2 + x_4^2 \leq 1 \end{array} \right] \\
 & \text{exactly}(1, Y) \\
 & 0 \leq x_1 \leq 8, \quad 0 \leq x_2 \leq 5.5 \\
 & -1 \leq x_3 \leq 1, \quad -1 \leq x_4 \leq 1 \\
 & Y_i \in \{\text{True}, \text{False}\}
 \end{aligned}$$



MIP Form (Big M)

$$\begin{aligned}
 & \min f(x_1, x_2) \\
 & \text{s.t.} \quad (x_1 - 1)^2 + (x_2 - 1)^2 \\
 & \quad \quad \quad + x_3^2 + x_4^2 \leq 1 + M_1(1 - Y_1) \\
 & \quad \quad \quad (x_1 - 4)^2 + (x_2 - 4.5)^2 \\
 & \quad \quad \quad + x_3^2 + x_4^2 \leq 1 + M_2(1 - Y_2) \\
 & \quad \quad \quad (x_1 - 7)^2 + (x_2 - 1)^2 \\
 & \quad \quad \quad + x_3^2 + x_4^2 \leq 1 + M_3(1 - Y_3) \\
 & 0 \leq x_1 \leq 8, \quad 0 \leq x_2 \leq 5.5 \\
 & -1 \leq x_3 \leq 1, \quad -1 \leq x_4 \leq 1 \\
 & Y_1 + Y_2 + Y_3 = 1, \quad Y_i \in \{0, 1\}
 \end{aligned}$$

REFORMULATION TECHNIQUES

GDP → MIP → Solve

- Some methods require **subproblems**

GDP → **Subproblems** → MIP → Solve

- Cutting Planes** solves for linear cuts NEW
- Multiple Big M** solves for coefficients NEW

Big-M (3 constants total: M_1, M_2, M_3)

$$r_1(z) \leq M_1 (1 - y_1)$$

$$r_2(z) \leq M_2 (1 - y_2)$$

$$r_3(z) \leq M_3 (1 - y_3)$$

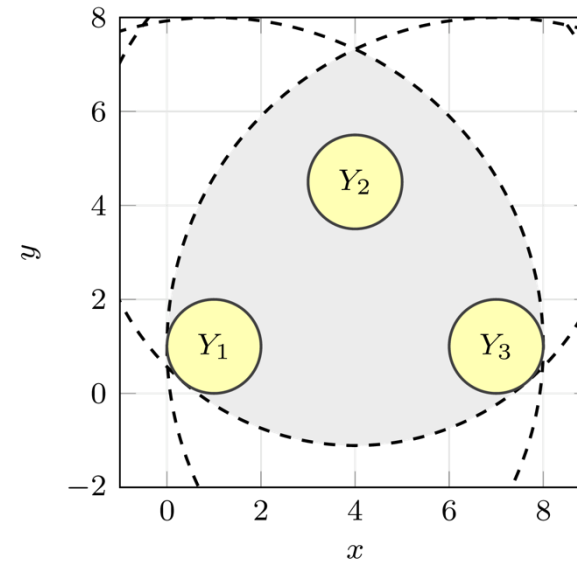
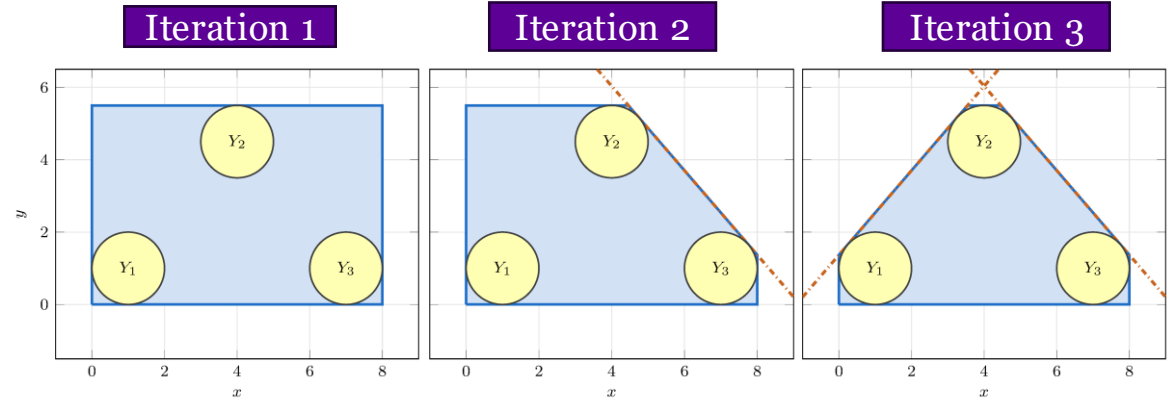
MBM (6 constants total: $M_{12}, M_{13}, M_{21}, M_{23}, M_{31}, M_{32}$)

$$r_1(z) \leq M_{12} y_2 + M_{13} y_3$$

$$r_2(z) \leq M_{21} y_1 + M_{23} y_3$$

$$r_3(z) \leq M_{31} y_1 + M_{32} y_2$$

Each M_{ij} is optimized for



DisjunctiveProgramming.jl

GDP Form

$\min_{x, Y} f(x_1, x_2)$

s.t.
$$\begin{aligned} & \left[\begin{array}{c} Y_1 \\ (x_1-1)^2 + (x_2-1)^2 + x_3^2 + x_4^2 \leq 1 \end{array} \right] \\ & \vee \\ & \left[\begin{array}{c} Y_2 \\ (x_1-4)^2 + (x_2-4.5)^2 + x_3^2 + x_4^2 \leq 1 \end{array} \right] \\ & \vee \\ & \left[\begin{array}{c} Y_3 \\ (x_1-7)^2 + (x_2-1)^2 + x_3^2 + x_4^2 \leq 1 \end{array} \right] \end{aligned}$$

exactly(1, Y)

$$0 \leq x_1 \leq 8, \quad 0 \leq x_2 \leq 5.5$$

$$-1 \leq x_3 \leq 1, \quad -1 \leq x_4 \leq 1$$

$$Y_i \in \{\text{True}, \text{False}\}$$

```
using DisjunctiveProgramming, Gurobi

model = GDPModel(Gurobi.Optimizer)

lb = [ 0.0, 0.0, -1.0, -1.0]
ub = [ 8.0, 5.5, 1.0, 1.0]

@variable(model, lb[k] <= x[k = 1:4] <= ub[k])

I = 1:3
@variable(model, Y[I], Logical)

c = [(1.0, 1.0, 0.0, 0.0),
      (4.0, 4.5, 0.0, 0.0),
      (7.0, 1.0, 0.0, 0.0)]

f(x1, x2) = (x1 - 0.0)^2 + (x2 - 5.0)^2

@constraint(model, [i = I],
            sum((x[k] - c[i][k])^2 for k = 1:4) <= 1,
            Disjunct(Y[i]))

@disjunction(model, Y)
@objective(model, Min, f(x[1], x[2]))

optimize!(model, gdp_method = Hull())
```

Supported Reformulations:

Big-M
Hull

Indicator

P-Split NEW

Cutting Planes NEW

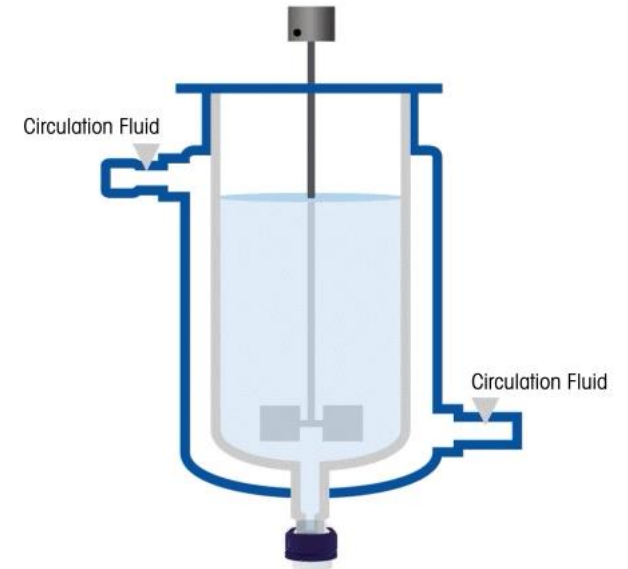
Multiple Big-M NEW

and their Infinite versions!

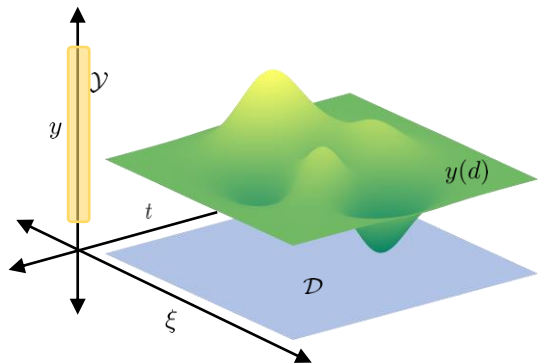


What is Infinite-dimensional Optimization (Infiniteopt) ?

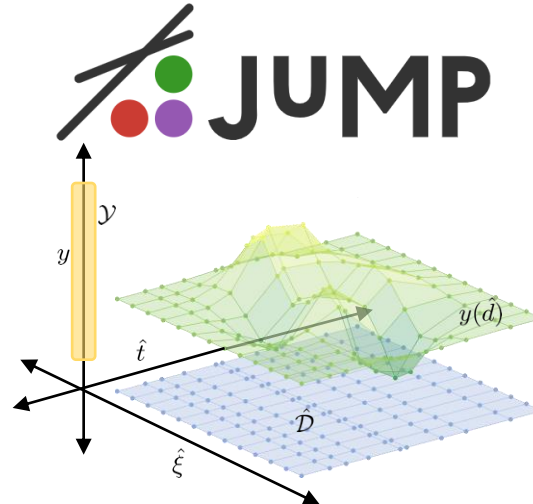
- Many real-world processes evolve over **continuous domains** like time and space
 - Chemical processes, rocket trajectory, portfolio optimization, etc...
 - Involves decision variables that are **infinite-dimensional functions**
 - **Example:** Optimal control of a jacketed reactor
- Implemented in `InfiniteOpt.jl`



 **InfiniteOpt**



Direct Transcription



Decision variables:

$$C_A(t), \quad T_{tank}(t), \quad T_{jacket}(t)$$



INFINITE DISJUNCTIVE PROGRAMS

- Modeling GDPs over infinite domains



InfiniteOpt

$$\bigvee_{i \in \mathcal{I}_j} \left[\begin{array}{c} W_{ij}(d) \\ r_{ij}(Dy, y(d), z, d) \leq 0 \end{array} \right], \quad d \in \mathcal{D}, j \in \mathcal{J}$$
$$\Omega(W(d)) = \text{True}, \quad d \in \mathcal{D}$$
$$W_{ij}(d) \in \{\text{True}, \text{False}\}, \quad d \in \mathcal{D}, i \in \mathcal{I}_j, j \in \mathcal{J}$$

- DisjunctiveProgramming.jl automates reformulation
- InfiniteOpt.jl automates transcription
- Techniques are adapted to the infinite setting

```
using DisjunctiveProgramming, InfiniteOpt, HiGHS

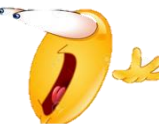
# Create the model

# Create the infinite variables

# Add the disjunctions and their indicator variables

# Add the logical propositions




# Reformulate and solve
```

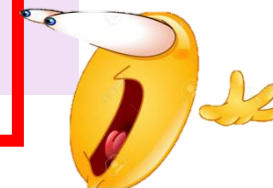


WHY INFINITEDISJUNCTIVEPROGRAMMING.JL?



- InfiniteDisjunctiveProgramming.jl adapts techniques for you
- Others like PyomoDAE do not bridge GDP and InfiniteOpt

	 PYOMO GDP	 G A M S	 <i>DisjunctiveProgramming.jl</i>
GDP Support	✓	✓	✓
InfiniteGDP	✗	✗	✓
Free	✓	✗	✓



CASE STUDY: SIMULATING SILVERSTONE CIRCUIT

Using `InfiniteDisjunctiveProgramming.jl`



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BACKGROUND

- 50/50 battery to fuel split in 2026
- **State of charge** is a real concern



4 Modes → 4 disjuncts over space/time

- **Regeneration:** Recharging battery through brake
- **Internal Combustion Engine (ICE):** Using fuel
- **Boost:** Electric + ICE
- **Clipping:** ICE + Charging Battery

Photo by Ben Sutherland CC BY-SA 2.0



MATH AND CODE

$$\min \int_0^L \frac{dt}{ds} ds$$

$$\text{s.t. } v(s) \frac{dt}{ds} = 1$$

$$m v(s) \frac{dv}{ds} = F_e(s) - \frac{1}{2} \rho C_d A v(s)^2 - f_r m g$$

$$E_b v(s) \frac{dSOC}{ds} = -P_b(s)$$

$$P_m(s) = F(s) v(s)$$

$$\bar{v}(\kappa) = \sqrt{\frac{\mu g}{\kappa - \frac{\mu \rho C_L A}{2m}}}$$

$$v_{\min} \leq v(s) \leq \min\{v_{\max}, \bar{v}(\kappa(s))\}$$

$$SOC_{\min} \leq SOC(s) \leq SOC_{\max}$$

$$\bigvee_{i \in \{\text{clip, ICE, boost, regen}\}} \begin{bmatrix} Y_i(s) \\ \underline{F}_i \leq F(s) \leq \bar{F}_i \\ \underline{P}_m^i \leq P_m(s) \leq \bar{P}_m^i \\ F_e(s) = \eta_{dt}^i F(s) \\ \underline{P}_b^i \leq P_b(s) \leq \bar{P}_b^i \\ g_i(P_m(s), P_b(s)) \leq 0 \end{bmatrix}$$

$$\text{EXACTLY}(1, Y_{\text{clip}}(s), Y_{\text{ice}}(s), Y_{\text{boost}}(s), Y_{\text{regen}}(s))$$

$$v(0) = v(L), \quad SOC(0) = SOC(L), \quad t(0) = 0$$

$$\forall s \in [0, L]$$

```

model = InfiniteGDPModel(Gurobi.Optimizer)

@infinite_parameter(model, s in [0, L], num_supports = N)
@objective(model, Min, integral(deriv(t, s), s))
@constraint(model, v * deriv(t, s) == 1)

@constraint(model, m_car * v * deriv(v, s) == F_eff - C_drag * v^2 -
    ↪ C_roll)
@constraint(model, E_batt * v * deriv(SOC, s) == -P_batt)
@constraint(model, P_motor == F_motor * v)

v_max(s) = sqrt(mu * g / (kappa(s) - mu * rho * ClA / (2 * m_car)))
@variable(model, v_min_global ≤ v ≤ v_max_global, Infinite(s))
@constraint(model, v ≤ v_max_pf)
@variable(model, SOC_min ≤ SOC ≤ SOC_max, Infinite(s))

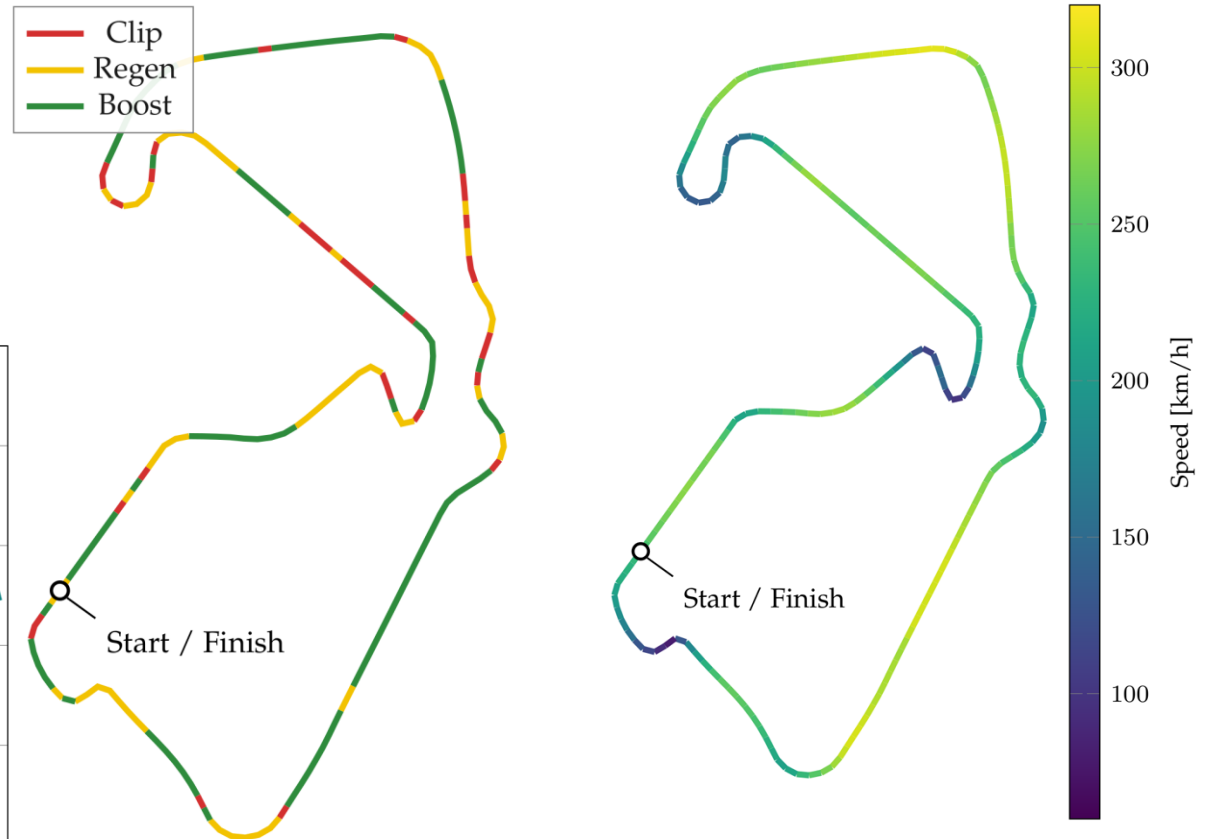
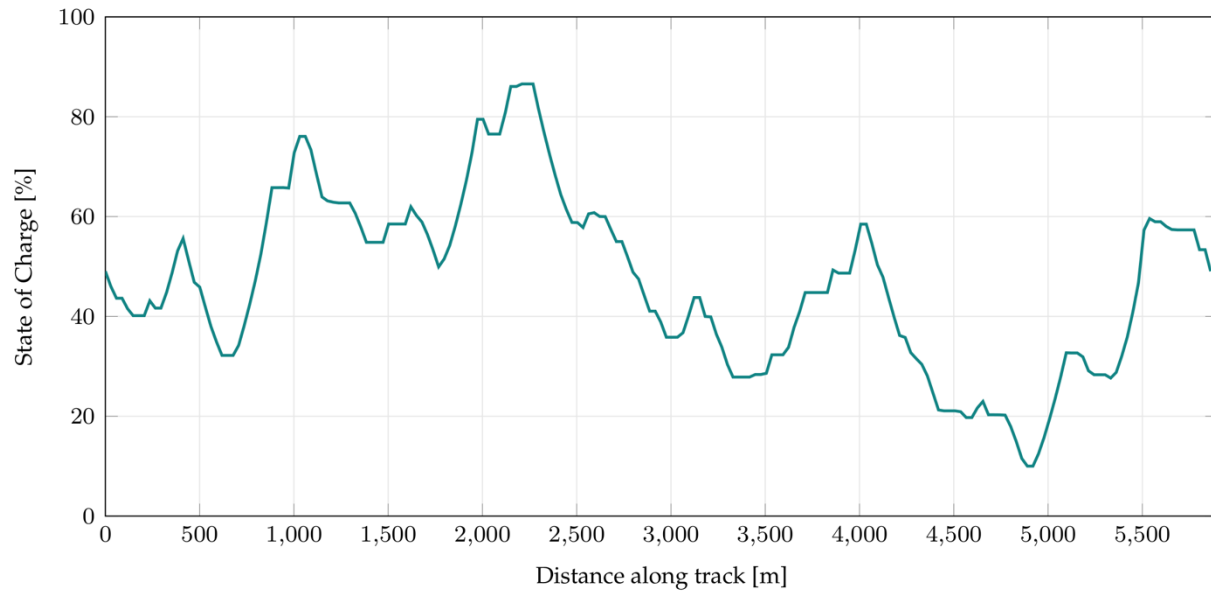
@variable(model, Y_clip, InfiniteLogical(s))
@variable(model, Y_ice, InfiniteLogical(s))
@variable(model, Y_boost, InfiniteLogical(s))
@variable(model, Y_regen, InfiniteLogical(s))
...
@constraint(model, F_motor ≥ F_motor_min_engine, Disjunct(Y_ice))
@constraint(model, P_motor ≤ P_ICE_max, Disjunct(Y_ice))
...
@disjunction(model, [Y_clip, Y_ice, Y_boost, Y_regen], exactly1 =
    ↪ true)
@constraint(model, v(0) == v(L))
@constraint(model, SOC(0) == SOC(L))
@constraint(model, t(0) == 0)

optimize!(model, gdp_method = MBM(Gurobi.Optimizer))

```

Results

- State of charge changes over time
- Mode selection is intuitively sound
- One decision (ICE) proved to be redundant



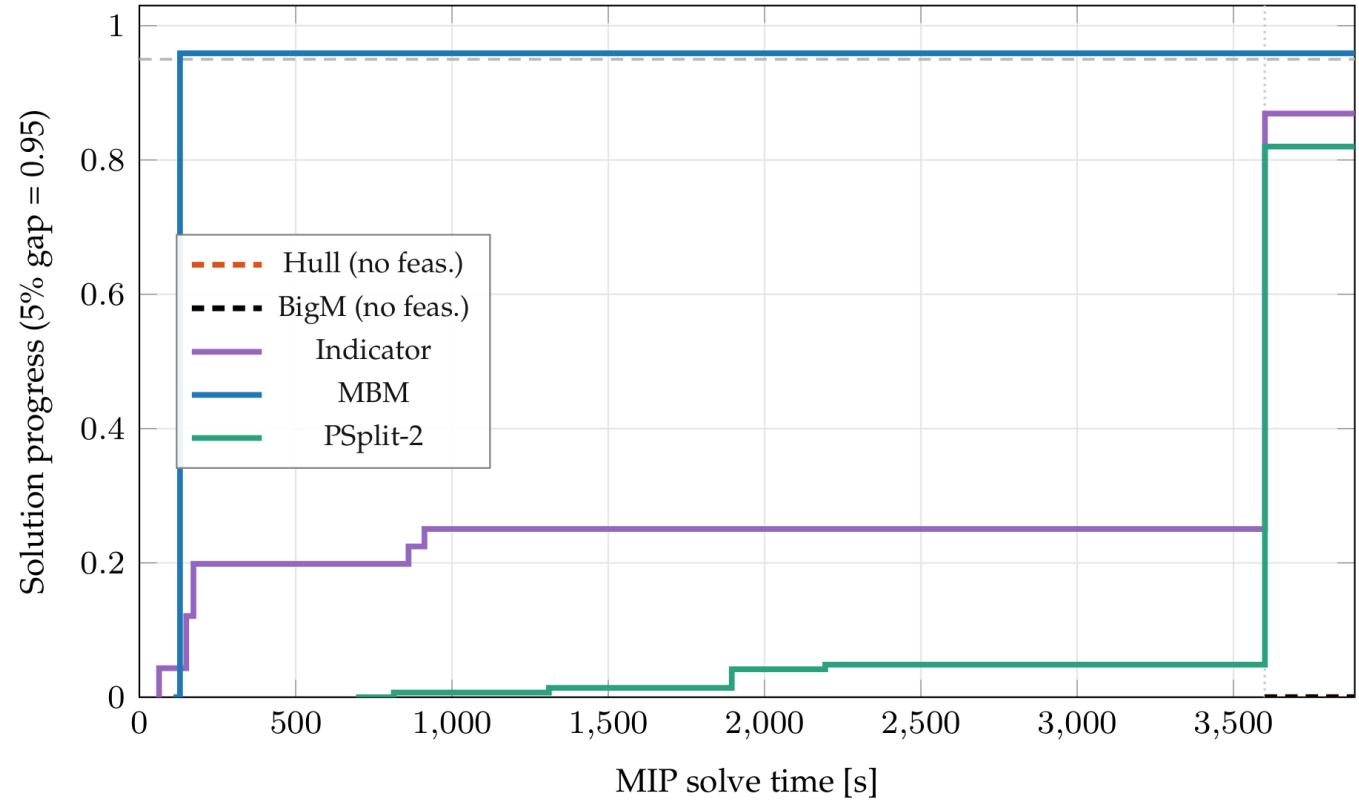
METHOD SOLVE TIMES

- **Hull:** No feasible solution found
- **Big-M:** No feasible solution found
- **Indicator:** 13.1% gap in 1 hour

NEW **Multiple Big-M:** Optimal in ~2 minutes

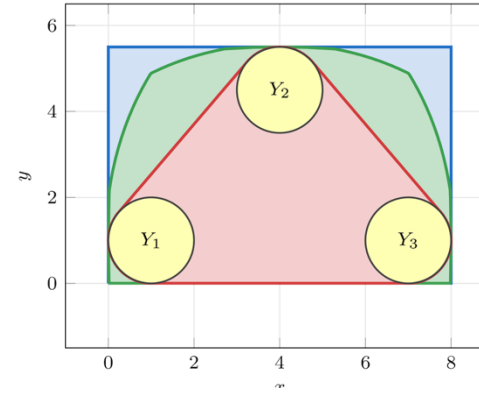
NEW **P-Split:** 18.0% gap in 1 hour

NEW **Cutting Planes:** Inappropriate technique

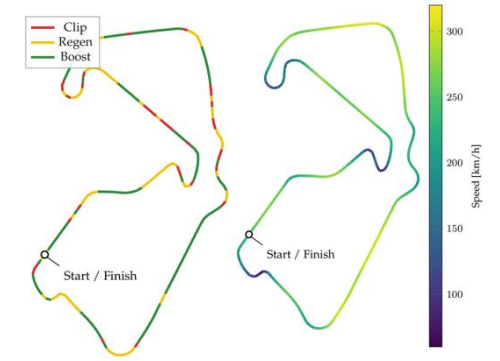
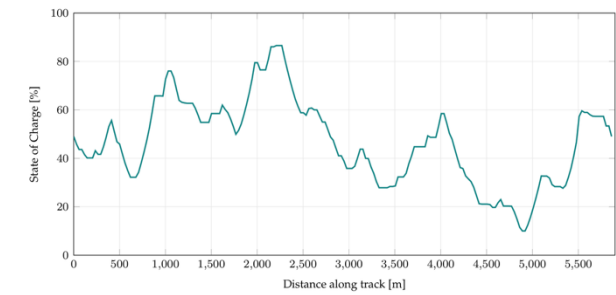


RECAP

- Big M and Hull are extremes of MIP relaxation
- Other methods exist to compromise
- Method switching is tedious
- DisjunctiveProgramming.jl automates GDP to MIP
- InfiniteDisjunctiveProgramming enables dynamic GDP



```
# --- 2. Pick a reformulation method ---  
optimize!(model, gdp_method = BigM())
```



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